

Oja Medians and Centers of Mass*

Dan Chen[†] Olivier Devillers[‡] John Iacono[§] Stefan Langerman[¶] Pat Morin^{*}

1 Introduction

Given a set S of n points in \mathbb{R}^d , the *Oja depth* [7] of a point $x \in \mathbb{R}^d$ is

$$d(x, S) = \sum_{y_1, \dots, y_d \in \binom{S}{d}} v(x, y_1, \dots, y_d) ,$$

where $v(p_1, \dots, p_{d+1})$ denotes the volume of the simplex whose vertices are $p_1 \dots p_{d+1}$.¹ A point in \mathbb{R}^d with the minimum Oja depth is called an *Oja center*.

1.1 New Results

In this paper we consider relationships between centers of mass of certain sets and Oja depth. The *center of mass* of a finite point set $S \subset \mathbb{R}^d$ is the average of those points,

$$c(S) = |S|^{-1} \sum_{x \in S} x .$$

If $P \subset \mathbb{R}^d$ is a bounded object of non-zero volume, the center of mass of P is

$$c(P) = \frac{\int_{x \in P} x \, dx}{v(P)} .$$

In this paper, we prove the following results about the Oja depth of an n point set S , whose convex hull A has unit volume and that has an Oja center x :

$$d(c(A), S) \leq \binom{n}{d} / (d + 1) , \tag{1}$$

$$d(c(S), S) \leq (d + 1) d(x, S) . \tag{2}$$

The bound in (1) is not known to be tight. The bound in (2) is tight, up to a lower-order term, for some point sets S .

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[†]School of Computer Science, Carleton University, dchen4@connect.carleton.ca, morin@scs.carleton.ca

[‡]INRIA Sophia Antipolis Méditerranée, France. Olivier.Devillers@sophia.inria.fr

[§]Department of Computer and Information Science, Polytechnic University, jiacono@poly.edu

[¶]Département d'Informatique, Université Libre de Bruxelles, stefan.langerman@ulb.ac.be

¹In Oja's original definition, the sum is normalized by dividing by $\binom{|S|}{d}$. We omit this here since it changes none of our results and clutters our formulas.

1.2 Related Results

Our first result, (1), is a form of *Centerpoint Theorem* that upper-bounds the Oja depth of $c(A)$, and hence also the Oja depth of x , in terms of the volume of the convex hull of S . Previously, centerpoint theorems were known for other depth functions such as Tukey depth [5, 8, 10] and simplicial depth [2, 3, 4]. To the best of our knowledge, this is the first such result for Oja depth.

Our next result, (2), can be viewed in two ways. The first is as a linear-time constant factor approximation for finding an Oja median.

In 1-d, Oja depth is minimized by the median, which can be found in $O(n)$ time. However, in 2-d, the best known algorithm for minimizing Oja depth exactly takes $O(n \log^3 n)$ time [1]. Approximation algorithms for minimizing Oja depth, based on uniform grids and sampling from $\binom{S}{d}$, are given by Ronkainen, Oja, and Orponen [9]. However, in pathological cases, their approximation algorithm is not guaranteed (or even likely) to find a point that closely approximates the Oja median, either in terms of distance or in terms of its Oja depth.²

Another view of (2) is that it gives insight into the Oja depth function and the Oja median. In some sense, it tells us that the Oja median is not terribly different from the center of mass of S , since the center of mass of S minimizes, to within a constant factor, the Oja depth function.

2 Oja Center and Mass Center of A

In this section, we relate the Oja depth of the center of mass of the convex hull of S to the volume of the convex hull of S . Throughout this section, A denotes the convex hull of S and we assume, without loss of generality, that $v(A) = 1$.

Our upper-bound is based on the following *central identity*: For any disjoint sets $X, Y \subseteq \mathbb{R}^d$ with $v(X \cup Y) > 0$,

$$c(X \cup Y) = \frac{v(X) c(X) + v(Y) c(Y)}{v(X \cup Y)} .$$

We first give an inductive proof of our result for point sets in \mathbb{R}^2 , and then give a proof for point sets in \mathbb{R}^d that uses tools from convex geometry.

²This follows from the fact that the value of the Oja depth function and the location of the Oja median can be arbitrarily different for two sets S_1 and S_2 that differ in only d points [6].

2.1 An Upper Bound in \mathbb{R}^2

Lemma 1 *Let E be a convex polygon, let g_e be the mass center of E , and let y_1, y_2 be any two points in E . Then $v(y_1 y_2 g_e) \leq \frac{1}{3} v(E)$.*

Proof. Assume, without loss of generality, that $y_1 y_2$ is horizontal and that $c(E)$ is above $y_1 y_2$. We may assume that $y_1 y_2$ is edge of E since, otherwise, we can remove the part of E that is below $y_1 y_2$ decreasing $v(E)$ and increasing $v(c(E) y_1 y_2)$.

The proof is by induction on the number of vertices of E . If E is a triangle then one can easily verify the result. Therefore, assume E has $n \geq 4$ vertices. Consider an edge ab of E where $a \neq y_2$ is adjacent to y_1 and let $c \neq a$ be adjacent to b (see Figure 1). Draw a ray r whose origin is at y_1 and such that the triangle t_1 supported by $y_1 a$, ab and r and the triangle t_2 supported by r , ab , and the line through bc have the same area. Such a ray is guaranteed to exist by a standard continuity argument that starts with r containing a and rotates about y_1 until r contains b .

Now, convert E into a polygon E' by removing t_1 and adding t_2 . This does not change the area of E . Furthermore, since t_1 and t_2 are separated by a vertical line, with t_2 above t_1 , this implies that the $c(E')$ has a larger y -coordinate than $c(E)$, so $v(c(E) y_1 y_2) \leq v(c(E') y_1 y_2)$. Note, also, that E' has $n - 1$ vertices so, by induction,

$$v(c(E) y_1 y_2) \leq v(c(E') y_1 y_2) \leq \frac{1}{3} v(E') = \frac{1}{3} v(E) ,$$

completing the proof. □

Theorem 2 *Let S be a set of points in \mathbb{R}^2 whose convex hull, A , has unit area. Then $d(c(A), S) \leq n^2/6$.*

Proof. According to Lemma 1,

$$d(c(A), S) = \sum_{y_1, y_2 \in \binom{S}{2}} v(y_1 y_2 c(A)) \leq \binom{n}{2} / 3 = \frac{n^2}{6} .$$

□

2.2 An Upper Bound in \mathbb{R}^d

First let us introduce a notion from convex geometry. Let A be a convex body in \mathbb{R}^d , where $d \geq 2$. Suppose A lies between parallel hyperplanes $x_1 = a$ and $x_1 = b$, where $a < b$. For each x with $a \leq x \leq b$, let A_x be the intersection of A with the hyperplane $x_1 = x$, and define r_x by the equation

$$\omega_{d-1} r_x^{d-1} = v_{d-1}(A_x),$$

where $v_{d-1}(X)$ denotes the $(d - 1)$ -dimensional volume of X and ω_{d-1} is the $(d - 1)$ -dimensional volume of the

unit $(d - 1)$ -ball. In this way, r_x is the radius of a $(d - 1)$ -ball whose v_{d-1} -volume is the same as that of A_x . For each $a \leq x \leq b$, let C_x be defined by the equation

$$C_x = \{(x, x_2, \dots, x_d) : x_2^2 + \dots + x_d^2 \leq r_x^2\}.$$

Then the set

$$C = \cup(C_x : a \leq x \leq b)$$

is called the *Schwarz rotation-symmetral* of A in the x_1 -axis. For example, in \mathbb{R}^3 , C is a stack of disks perpendicular to, and centered on, the x -axis. Each disk has the same area as the corresponding slice of A .

Theorem 3 (Webster [11]) *Let A be a convex body in \mathbb{R}^d ($d \geq 2$) whose Schwarz rotation-symmetral in the x_1 -axis is C . Then C is a convex body having the same volume as A .*

Lemma 4 *Let g be the center of mass of a convex d -polytope P . Then any d -simplex T whose vertices are g plus d points inside P has volume at most $v(P)/(d + 1)$.*

Proof. Let p_1, \dots, p_d be any d points in P , and let h be the hyperplane that contains them. If g is in h , then $v(T) = 0$. If not, rotate P to make h perpendicular to the x_1 -axis with g above h . If there is any volume of P below h , we can cut that part off from P to obtain a new polytope P' . The volume of P' will be less than 1, and its mass center g' will be above g . In this way, if (P, p_1, \dots, p_d) is a counterexample to the lemma, then so is (P', p_1, \dots, p_d) . The face of P' in h is a convex hull B containing the d points. Let q be a point above h such that the pyramid D with B as base and q as apex has the same volume as P' .

Let C be the Schwarz rotation-symmetral of P' in the x_1 -axis, and R be that of D (see Figure 2). Note that R is a conic pyramid. By Theorem 3, C is convex and $v(C) = v(P') = v(R)$. Let c be the intersection of the surfaces of C and R above B . Note that the surface of R is bounded by a collection of lines that pass through q . Each of these lines intersects C in at most 2 points. One of these points has the same x_1 -coordinate as B and the other points lie on the boundary of a $(d - 1)$ -ball c .

Let x_c be the x_1 -coordinate of c . Since C is convex, the surface of C below $x_1 = x_c$ is outside the surface of R . By the definition of Schwarz rotation-symmetral, the volume of C that is outside of R is below x_c , and the volume of R outside of C is above x_c . Therefore, the mass center of R is above that of C because of central identity.

The mass centers of P' and C have the same height because in the Schwarz rotation-symmetral C_x has the same x_1 value as A_x . So do the mass centers of D and R . Let g_d be mass center of D . Since D is a pyramid,

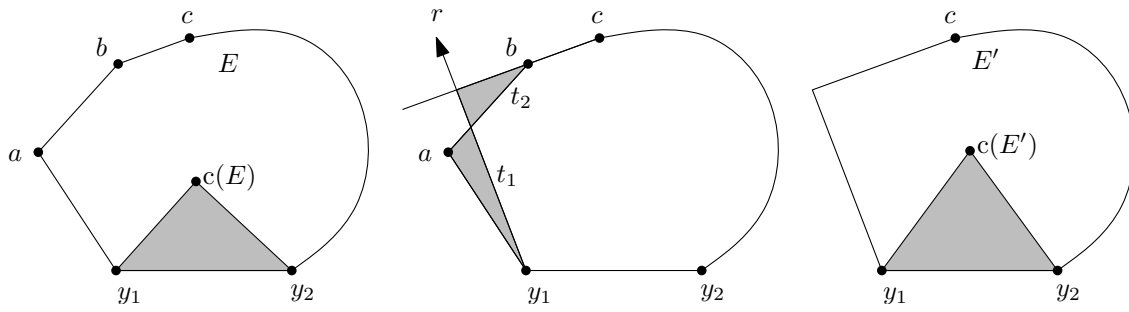


Figure 1: The proof of Lemma 1.

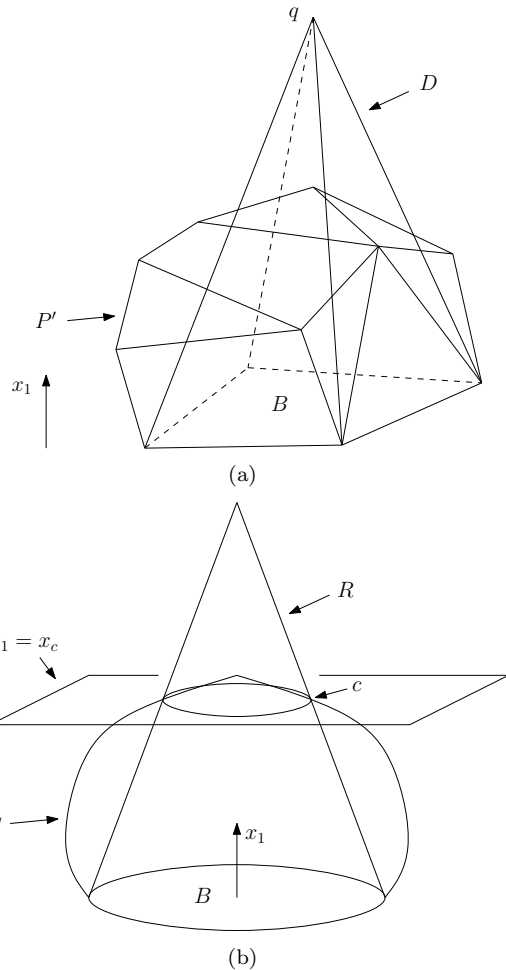


Figure 2: The Schwarz rotation-symmetral of P' and D

the convex hull of the d points is contained in B , g is below g' , and g' is below g_d , we have

$$v(T) \leq v(P')/(d+1) \leq v(P)/(d+1).$$

To see this, consider that $v(T) = v(g, p_1, \dots, p_d) \leq v(g_d, p_1, \dots, p_d) = v(q, p_1, \dots, p_d)/(d+1)$. \square

The bound in Lemma 4 is tight, for example, when S consists of the $d+1$ vertices of a simplex. Next we

show how this relates to Oja depth:

Theorem 5 *Let S be a set of points in \mathbb{R}^d whose convex hull, A , has unit volume. Then $d(c(A), S) \leq \binom{n}{d}/(d+1)$.*

Proof.

$$\begin{aligned} d(c(A), S) &= \sum_{y_1, \dots, y_d \in \binom{S}{d}} v(c(A), y_1, \dots, y_d) \\ &\leq \binom{n}{d}/(d+1), \end{aligned}$$

where the inequality is an application of Lemma 4. \square

3 Oja Center and Mass Center of S

In this section, we show that the center of mass of S provides a constant-factor approximation to the point of minimum Oja depth.

Theorem 6 *For any finite set $S \subset \mathbb{R}^d$, and any $x \in \mathbb{R}^d$, $d(c(S), S) \leq 2d(x, S)$.*

Proof. Denote the elements of S by p_1, \dots, p_n in any order. Let the multiset S_i contain p_1, \dots, p_i as well as $n-i$ copies of x . Let $c_i = c(S_i)$. We will show, by induction on i , that $d(c_i, S_i) \leq 2d(x, S_i)$ for all $i \in \{0, \dots, n\}$. This is sufficient, since $S_n = S$.

For the base case S_0 consists of n copies of x , so $c_0 = x$ and $d(c_0, S_0) = 0 = 2d(x, S_0)$. Next, we assume that $d(c_i, S_i) \leq 2d(x, S_i)$ and prove that $d(c_{i+1}, S_{i+1}) \leq 2d(x, S_{i+1})$. Note that

$$d(x, S_{i+1}) = d(x, S_i) + |p_{i+1} - x|.$$

Furthermore,

$$c_{i+1} = c_i + (p_{i+1} - x)/n,$$

so

$$\begin{aligned}
 d(c_{i+1}, S_{i+1}) &= d(c_i, S_i) \\
 &\quad + \sum_{q \in S_i} (|c_{i+1} - q| - |c_i - q|) \\
 &\quad + (|c_{i+1} - p_{i+1}| - |c_{i+1} - x|) \\
 &\leq d(c_i, S_i) + n|p_{i+1} - x|/n \\
 &\quad + (|c_{i+1} - p_{i+1}| - |c_{i+1} - x|) \\
 &\leq d(c_i, S_i) + 2|p_{i+1} - x| \\
 &\leq 2d(x, S_i) + 2|p_{i+1} - x| \\
 &= 2d(x, S_{i+1}) ,
 \end{aligned}$$

as required. \square

We remark that the above proof uses little more than triangle inequality. In particular, the same proof shows that the center of mass gives a 2-approximation for the Fermat-Weber center in any dimension.³ Unfortunately, in higher dimensions, Oja depth does not enjoy this nice property.

Theorem 7 For any finite set $S \subseteq \mathbb{R}^d$, $d(c(S), S) \leq (d + 1)d(x, S)$ for any $x \in \mathbb{R}^d$.

Proof. In this proof, we will make use of the fact that, for any d -simplex T with vertex set V_T and a point $q \in \mathbb{R}^d$,

$$v(T) \leq \sum_{p_1, \dots, p_d \in \binom{V_T}{d}} v(p_1, \dots, p_d, q) , \quad (3)$$

since T is contained in the union of the simplices on the right hand side. Equality occurs if q is inside T .

Define S_i as in the proof of Theorem 6. Let S' be S_{i+1} with one occurrence of p_{i+1} removed. The induction and base case are the same as in Theorem 6. First, we have

$$\begin{aligned}
 d(x, S_{i+1}) &= d(x, S_i) \\
 &\quad + \sum_{Q \in \binom{S_i}{d-1}} v(x, p_{i+1}, Q), \quad (4)
 \end{aligned}$$

where Q is a set of $d - 1$ points, and

$$\begin{aligned}
 d(c_{i+1}, S_{i+1}) &= d(c_i, S_i) \\
 &\quad + \sum_{P \in \binom{S_i}{d}} (v(c_{i+1}, P) - v(c_i, P)) \quad (5)
 \end{aligned}$$

$$\begin{aligned}
 &\quad + \sum_{Q \in \binom{S'_i}{d-1}} (v(c_{i+1}, p_{i+1}, Q) - v(c_{i+1}, x, Q)), \quad (6)
 \end{aligned}$$

where P is a set of d points. We denote y^\perp the projection of a point y on a line perpendicular to the $d - 1$ dimensional simplex P .

$$\begin{aligned}
 &|v(c_{i+1}, P) - v(c_i, P)| \\
 &= (1/d)v_{d-1}(P) \left| \|c_i^\perp P^\perp\| - \|c_{i+1}^\perp P^\perp\| \right| \\
 &\leq (1/d)v_{d-1}(P) \|c_i^\perp - c_{i+1}^\perp\| \\
 &\leq (1/d)v_{d-1}(P) \|(1/n)x^\perp p_{i+1}^\perp\|
 \end{aligned}$$

Then if x^\perp and p_{i+1}^\perp are on the same side of the hyperplane supporting P , we have

$$\begin{aligned}
 &(1/d)v_{d-1}(P) \|(1/n)x^\perp p_{i+1}^\perp\| \\
 &\leq (1/nd)v_{d-1}(P) \left| \|x^\perp P^\perp\| - \|p_{i+1}^\perp P^\perp\| \right| \\
 &\leq (1/n) |v(p_{i+1}, P) - v(x, P)| \\
 &\leq (1/n) \sum_{Q \in \binom{P}{d-1}} v(x, p_{i+1}, Q)
 \end{aligned}$$

Otherwise if x^\perp and p_{i+1}^\perp are on different sides of the hyperplane supporting P , we have $\|x^\perp p_{i+1}^\perp\| = \|x^\perp P^\perp\| + \|p_{i+1}^\perp P^\perp\|$. In this case the two simplices Px and Pp_{i+1} are disjoint and the convex hull of Pxp_{i+1} is covered by the union of the simplices Qxp_{i+1} for $Q \in \binom{P}{d-1}$, thus

$$\begin{aligned}
 &(1/d)v_{d-1}(P) \|(1/n)x^\perp p_{i+1}^\perp\| \\
 &\leq (1/nd)v_{d-1}(P) (\|x^\perp P^\perp\| + \|p_{i+1}^\perp P^\perp\|) \\
 &\leq (1/n) v(p_{i+1}, P) + v(x, P) \\
 &\leq (1/n) \sum_{Q \in \binom{P}{d-1}} v(x, p_{i+1}, Q)
 \end{aligned}$$

We can now prove that (5) \leq (4) as follows:

$$\begin{aligned}
 &\sum_{P \in \binom{S_i}{d}} (v(c_{i+1}, P) - v(c_i, P)) \\
 &\leq \sum_{P \in \binom{S_i}{d}} (|v(c_{i+1}, P) - v(c_i, P)|) \\
 &\leq (1/n) \sum_{P \in \binom{S_i}{d}} \sum_{Q \in \binom{P}{d-1}} v(x, p_{i+1}, Q) \\
 &\leq (n - (d - 1)/n) \sum_{Q \in \binom{S_i}{d-1}} v(x, p_{i+1}, Q) .
 \end{aligned}$$

³The Fermat-Weber center of a point set S in \mathbb{R}^d is the point x that minimizes $\sum_{y \in S} \|x - y\|$.

Next, we show that $(6) \leq d \times (4)$. Applying (3),

$$\begin{aligned} & \sum_{Q \in \binom{S'}{d-1}} (v(c_{i+1}, p_{i+1}, Q) - v(c_{i+1}, x, Q)) \\ & \leq \sum_{Q \in \binom{S'}{d-1}} \left(v(x, p_{i+1}, Q) + \sum_{R \in \binom{Q}{d-2}} v(x, p_{i+1}, c_{i+1}, R) \right) \\ & \leq \sum_{Q \in \binom{S_i}{d-1}} v(x, p_{i+1}, Q) \\ & \quad + (n - 1 - (d - 2)) \sum_{R \in \binom{S_i}{d-2}} v(x, p_{i+1}, c_{i+1}, R) , \end{aligned}$$

where R is a set of $d - 2$ points. By linearity of the determinant we have

$$\begin{aligned} v(x, p_{i+1}, c_{i+1}, R) &= (1/n) \sum_{y \in S_{i+1}} v(x, p_{i+1}, y, R) \\ &= (1/n) \sum_{y \in S_i} v(x, p_{i+1}, y, R) \end{aligned}$$

Since the absolute value of the sum can be bounded by the sum of the absolute values, we get

$$v(x, p_{i+1}, c_{i+1}, R) \leq (1/n) \sum_{y \in S_i} v(x, p_{i+1}, y, R),$$

and thus

$$\sum_{R \in \binom{S_i}{d-2}} v(x, p_{i+1}, c_{i+1}, R) = (d-1/n) \sum_{Q \in \binom{S_i}{d-1}} v(x, p_{i+1}, Q).$$

Thus we can get $(6) \leq d \times (4)$.

Finally, we resubstitute to obtain

$$\begin{aligned} & d(c_{i+1}, S_{i+1}) \\ & \leq d(c_i, S_i) + (d + 1) \sum_{Q \in \binom{S_i}{d-1}} v(x, p_{i+1}, Q) \\ & \leq (d + 1) d(x, S_i) + (d + 1) \sum_{Q \in \binom{S_i}{d-1}} v(x, p_{i+1}, Q) \\ & = (d + 1) d(x, S_{i+1}) . \end{aligned}$$

Then we have $d(c(S), S) \leq (d + 1) d(x, S)$. □

We remark that Theorem 6 and 7 are essentially the best possible. To see this, take the multiset S that contains $n - d$ copies of the origin o , and each of the remaining d points has one different coordinate 1 and all other coordinates 0. In this case $d(o, S) = 1/d!$ and $d(c(S), S) = (d + 1 - O(d^2/n)) \times 1/d!$.

4 Conclusion

We have given several results on Oja depth and centers of mass. There are several directions for future work.

Theorem 5 has no matching lower bound. The best lower-bound we know is that placing $n/(d + 1)$ points at each vertex of any d -simplex of unit volume yields to an Oja depth of $n^d/(d + 1)^d$ for any point inside the simplex. For $d = 2$, for example, Theorem 5 implies $d(x, S) \leq n^2/6 - O(n)$ where as the best lower bound (above) has $d(x, S) \geq n^2/9$. This construction leads us to the following conjecture:

Conjecture 1 *For any point set $S \subset \mathbb{R}^d$ whose convex hull has unit volume, there exists $x \in \mathbb{R}^d$, such that $d(x, S) \leq n^d/(d + 1)^d$*

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