Review & Plan

COMP 7720 - Online Algorithms  Paging and $k$-Server Problem

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Today’s objectives

- Caching problem & advice
- $k$-server problem
  - Lower bound for deterministic algorithms
  - Greedy algorithms
  - Paths & trees
Caching Problem
There are two types of memory: a fast ‘cache’ of size $k$, and a slow memory of unbounded size.

- The input is an online sequence of requests to pages of size 1.

- To serve a request to page $x$, it should be in the cache
  - In case $x$ is not in the cache, a fault of cost 1 happens
  - In case $x$ is in the cache, a hit of cost 0 happens
  - The goal is to minimize the total number of faults

- To bring $x$ to the cache, we might need to evict a page.
  - A caching algorithm is defined through its eviction policy.
Caching problem: a review

- Latest-In-Future (LIF) is the optimal offline algorithm.
- No deterministic algorithm has a competitive ratio better than $k$. 
Caching Problem

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- An algorithm is marking if it maintains a ‘mark’ for each page.
  - After a request to $x$ mark it.
  - Always evict an unmarked page (if all marked, first unmark all pages and then evict one)
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  - After a request to $x$ mark it.
  - Always evict an unmarked page (if all marked, first unmark all pages and then evict one)
- Any deterministic marking algorithm has a competitive ratio of $k$.
  - Least-Recently-Used (LRU), and Flash-When-Full (FWF) both have competitive ratio $k$. 
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  - Always evict an unmarked page (if all marked, first unmark all pages and then evict one)
- Any deterministic marking algorithm has a competitive ratio of $k$.
  - Least-Recently-Used (LRU), and Flash-When-Full (FWF) both have competitive ratio $k$.
- Fist-In-First-Out (FIFO) also has a competitive ratio of $k$. 
A randomized algorithm which randomly evict a page has a competitive ratio of $k$.

A marking algorithm that evicts an unmarked page uniformly at random has a competitive ratio of $H_k$

- $H_k = 1 + 1/2 + 1/3 + \ldots + 1/k$
- For large values of $k$, we have $H_k \approx \ln(k) \in \Theta(\log k)$.

In fact, no randomized algorithm can achieve a better competitive ratio (i.e., $o(\log k)$).
Caching Problem

Caching problem: a review (cntd.)

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How many bits of advice are sufficient to achieve an optimal algorithm?

- $n$: the length of input sequence (number of requests).
- $k$: the size of the cache
- Hint: an algorithm has to make at most $n$ decisions about the page to be evicted.
  - one decision per fault.

At most $O(n \log k)$ bits are sufficient!
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What does the advice encode?

What is the size of advice?

How the algorithm works, provided by these bits of advice?

Why the algorithm has a competitive ratio of 1 (optimal here)?
Caching Problem

Caching & advice

- What does the advice encode? The advice indicates, for each request, what an optimal offline algorithm Opt evicts in case of a failure.

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- How the algorithm works, provided by these bits of advice?

- Why the algorithm has a competitive ratio of 1 (optimal here)?
What does the advice encode? The advice indicates, for each request, what an optimal offline algorithm \( \text{Opt} \) evicts in case of a failure.

What is the size of advice? Assume \( \text{Opt} \) makes \( m \leq n \) faults for the optimal algorithm. For each fault, the advice indicates which page should be evicted. There are \( k \) pages in the cache, and the evicted page can be indicated in \( \Theta(\log k) \) bits. The total number of bits will be \( m \cdot O(\log k) \in O(n \log k) \).

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How the algorithm works, provided by these bits of advice? It just mimics Opt; whenever there is a fault, it reads the advice to see which page should be evicted.

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How the algorithm works, provided by these bits of advice? It just mimics Opt; whenever there is a fault, it reads the advice to see which page should be evicted.

Why the algorithm has a competitive ratio of 1 (optimal here)? It works exactly like Opt.
Theorem

There is an online algorithm that, provided with $O(n \log k)$ bits of advice, can achieve an optimal solution.

- It is a naive solution :’-)
- Can we achieve an optimal solution with a smaller number of bits of advice?
  - For many problems, the answer is no!
  - For caching problem, we can indeed do better.
Caching Problem

Caching & advice (cntd.)

- Assume $\text{OPT}$ brings a page $x$ to the cache at time $t$.
  - Either $\text{OPT}$ evicts $x$ before the next access to $x \rightarrow x$ is mortal.
  - $\text{OPT}$ keeps $x$ in the cache until the next access to $x \rightarrow x$ is resident.
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If $\text{OPT}$ has a hit for request $x \Rightarrow x$ has been resident in cache since its last access to $x$. 

Consider an algorithm $\text{ResMor}$ that evicts a mortal page if an eviction is required. $\text{ResMor}$ always has the same resident pages as $\text{OPT}$ in its cache. The mortal pages might be different. $\text{OPT}$ and $\text{ResMor}$ have the same cost. Assume $\text{OPT}$ has smaller cost $\Rightarrow$ there is a request to $x$ that is a hit by $\text{OPT}$ and a miss for $\text{ResMor}$ $\Rightarrow x$ is resident in $\text{OPT}$ and not in $\text{ResMore}$ $\Rightarrow$ they maintain different resident pages (ResMore has evicted a resident page at some point) $\Rightarrow$ contradiction.
Caching Problem

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Assume with each request, there is one bit of advice that indicates whether the requested page is resident or mortal after the request.
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We can think of **ResMor** as an online algorithm with \( n \) bits of advice.
What does the advice encode?

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How the algorithm works, provided by these bits of advice?

Why the algorithm has a competitive ratio of 1 (optimal here)?
What does the advice encode? For each request there is one bit of advice indicating whether the requested page is evicted before the next request to it (i.e., it is mortal) or not (i.e., it is resident).

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How the algorithm works, provided by these bits of advice? When eviction is required, it evicts any mortal page.

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How the algorithm works, provided by these bits of advice? When eviction is required, it evicts any mortal page.

Why the algorithm has a competitive ratio of 1 (optimal here)? It maintains the same resident pages as Opt; so in case of a hit by Opt there will be a hit by the algorithm.
Advice complexity of paging

- With $n$ bits of advice, one can achieve an optimal algorithm.
- With roughly $\log\left(\frac{r+1}{r^{r+1}}\right) \cdot n$ bits, one can achieve a competitive ratio of $n$.
  - With roughly 0.27$n$ bits, one can achieve a competitive ratio of 2.
  - With roughly 0.24$n$ bits, one can achieve a competitive ratio of 3.
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- For a potential project, do a survey on advice complexity of paging, and try to deduce new results!
k-Server Problem
Introduction

\textbf{\textit{k-sever problem}}

- A metric is a set of points with a distance between each of pairs so that $d(x, y) \leq d(x, z) + d(z, y)$.

- E.g., a connected, undirected graph or a set of points in plane

\[
\sigma = \langle S, M, K, A, D, B, D, B, D \rangle
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\text{costs} = 2, 0, 2, 1, 1, 1, 1, 1, 1
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- We have a metric space of size \( m \)
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- **COMP 7720 - Online Algorithms**
- **Paging and k-Server Problem**
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Minimize the total distance moved by servers
The $k$-server problem

- What happens if we have a complete graph (a uniform metric)?
The $k$-server problem

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  - If there is a request to a vertex at which a server is located → there is no cost; otherwise, there is a cost of 1 to move a server to requested vertex.
Introduction

The k-server problem

- What happens if we have a complete graph (a uniform metric)?
  - If there is a request to a vertex at which a server is located → there is no cost; otherwise, there is a cost of 1 to move a server to requested vertex.
    - Think of vertices as pages; vertices with servers on them are pages in the cache → caching problem.
The \( k \)-server problem

What happens if we have a complete graph (a \textit{uniform} metric)?

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Recall that for caching problem, we have:

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\textbf{Theorem}
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No deterministic algorithm can achieve a competitive ratio better than \( k \), and LRU and FIFO achieve this ratio.
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- Recall that for caching problem, we have:

**Theorem**

No deterministic algorithm can achieve a competitive ratio better than $k$, and LRU and FIFO achieve this ratio.

No randomized algorithm can achieve a competitive ratio that is asymptotically better than $\Theta(\log k)$ and Mark algorithm achieves this.

- $k$-server problem has the right level of difficulty compared to paging (which is ‘too easy’) and Metrical Task Systems (another problem which is ‘too hard’).
Greedy Algorithm

Move the closest server to serve each request.
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Greedy Algorithm

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\[ \text{costs} = 2 \quad 0 \quad 2 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \quad 1 \]
Greedy Algorithm

- Move the closest server to serve each request.

\[
\sigma = < S, M, K, A, D, B, D, B, D >
\]

Costs:

\[
\text{costs} = 2, 0, 2, 1, 1, 1, 1, 1, 1
\]
Greedy Algorithm

- Move the closest server to serve each request.
- Is Greedy a good algorithm?
Greedy Algorithm

- Move the closest server to serve each request.
- Is Greedy a good algorithm?
  - what about the input \( \sigma = \langle B \ R \ B \ R \ldots \rangle \)?
    - For \( n \) requests, greedy incurs a cost of \( n \)
    - \( \text{OPT} \) moves another server from \( M \) to \( T \) at a cost of 3 and incurs no cost.
    - Competitive ratio will be at least \( \frac{n}{3} \) for this graph!
Introduction

Greedy Algorithm

Theorem

For any graph of diameter \( d \), the competitive ratio of greedy is at least \( \frac{n}{2d} \).

- It holds for any graph, even a path!
- Consider two vertices \( A \) and \( B \) which are close to one server and further from other servers.
  - Greedy servers sequence \( \langle A \ B \ A \ B \ldots \rangle \) by one server
Greedy Algorithm

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  - Greedy servers sequence \( \langle A \ B \ A \ B \ldots \rangle \) by one server
  - The cost of Greedy is at least \( n \)
  - An optimal algorithm moves two servers to \( A \) and \( B \) at a cost of at most \( 2d \).
  - The competitive ratio of greedy becomes at least \( \frac{n}{2d} \).
Competitive analysis

For general metrics

No deterministic online algorithm can be better than $k$-competitive. (we see the proof in the next class)
**Conjecture**

*Conjecture: for any metric space, there is a deterministic algorithm with competitive ratio of $k$.*

- $k$-server conjecture is one of the big open problems in the context of online algorithms.
Conjecture

Conjecture: for any metric space, there is a deterministic algorithm with competitive ratio of $k$.

- $k$-server conjecture is one of the big open problems in the context of online algorithms.
  - Verified when $k = 2$, $m = k + 1$, $m = k + 2$, and trees.
- In the next class, we learn about potential function algorithm, which has a competitive ratio of $2k - 1$